

Bi-Intermediate Logics of Trees and Co-trees

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Bi-intuitionistic logic bi-IPC is the conservative extension of intuitionistic logic IPC obtained by adding a new binary connective \leftarrow to the language, called the *co-implication* (or exclusion, or subtraction), which behaves dually to \rightarrow . In this way, bi-IPC achieves a symmetry, which IPC lacks, between the connectives $\wedge, \top, \rightarrow$ and \vee, \perp, \leftarrow , respectively.

The Kripke semantics of bi-IPC [25] provides a transparent interpretation of co-implication: given a Kripke model \mathfrak{M} , a point x in \mathfrak{M} , and formulas ϕ, ψ , then

$$\mathfrak{M}, x \models \phi \leftarrow \psi \iff \exists y \leq x (\mathfrak{M}, y \models \phi \text{ and } \mathfrak{M}, y \not\models \psi).$$

Equipped with this new connective, bi-IPC achieves significantly greater expressivity than IPC. For instance, if the points of a Kripke frame are interpreted as states in time, the language of bi-IPC is expressive enough to talk about the past, something that is not possible in IPC. With this example in mind, Wolter extended Gödel's interpretation of IPC into **S4** to an interpretation of bi-IPC into tense-S4 [30]. In particular, he proved a version of the Blok-Esakia Theorem [6, 13] stating that the lattice $\Lambda(\text{bi-IPC})$ *bi-intermediate logics* (i.e., consistent axiomatic¹ extensions of bi-IPC) is isomorphic to that of consistent normal tense logics containing **Grz.t**, see also [9, 28].

The greater symmetry of bi-IPC with respect to IPC is reflected in the fact that bi-IPC is algebraized in the sense of [7] by the variety bi-HA of *bi-Heyting algebras* [24], i.e., Heyting algebras whose order duals are also Heyting algebras. As a consequence, the lattice $\Lambda(\text{bi-IPC})$ is dually isomorphic to that of nontrivial varieties of bi-Heyting algebras. The latter, in turn, is amenable to the methods of universal algebra and duality theory because the category of bi-Heyting algebras is dually isomorphic to that of *bi-Esakia spaces* [12], see also [3].

The theory of bi-Heyting algebras was developed in a series of papers by Rauszer and others motivated by the connection with bi-intuitionistic logic. However, bi-Heyting algebras arise naturally in other fields of research as well such as topos theory [20, 21, 26]. Furthermore, the lattice of open sets of an Alexandrov space is always a bi-Heyting algebra, and so is the lattice of subgraphs of an arbitrary graph (see, e.g., [29]) and, similarly, every quantum system can be associated with a complete bi-Heyting algebra [11].

The lattice $\Lambda(\text{IPC})$ of intermediate logics (i.e., consistent extensions of IPC) has been thoroughly investigated (see, e.g., [8]). On the other hand, the lattice $\Lambda(\text{bi-IPC})$ of bi-intermediate logics lacks such an in-depth analysis, but for some recent developments see, e.g., [1, 4, 14, 15, 27]. In this paper we shall contribute to fill this gap by studying a simpler, yet nontrivial, sublattice of $\Lambda(\text{bi-IPC})$: the lattice of consistent extensions of the *bi-intuitionistic linear calculus* (or the bi-Gödel-Dummett's logic),

$$\text{bi-LC} := \text{bi-IPC} + (p \rightarrow q) \vee (q \rightarrow p).$$

¹From now on we will use *extension* as a synonym of *axiomatic extension*.

Notably, the properties of $\Lambda(\text{bi-IPC})$ and its extensions diverge significantly from those of its intermediate counterpart, i.e., the *intuitionistic linear calculus* (or the Gödel-Dummett's logic) $\text{LC} := \text{IPC} + (p \rightarrow q) \vee (q \rightarrow p)$ [10, 16].

The choice of bi-LC as a case study was motivated by some of its properties that make it an interesting logic on its own. On the one hand, bi-LC is complete in the sense of Kripke semantics with respect to the class of *co-trees* (i.e., order duals of trees). Moreover, we prove that the bi-intuitionistic logic of linearly ordered Kripke frames is a proper extension of bi-LC. This contrasts with the case of intermediate logics, where LC is both the logic of the class of linearly ordered Kripke frames and of co-trees. Because of this, the language of bi-IPC seems more appropriate to study tree-like structures than that of IPC. Furthermore, because of the symmetric nature of bi-intuitionistic logic, our results on extensions of bi-LC can be extended in a straightforward manner to the extensions of the bi-intermediate logic of trees by replacing in what follows every formula φ by its dual $\neg\varphi^\partial$, where φ^∂ is the formula obtained from φ by replacing each occurrence of $\wedge, \top, \rightarrow$ by \vee, \perp, \leftarrow respectively, and every algebra or Kripke frame by its order dual.

On the other hand, the logic bi-LC admits a form of a classical reductio ad absurdum, as we proceed to explain. A deductive system \vdash is said to have a *classical inconsistency lemma* if, for every nonnegative integer n , there exists a finite set of formulas $\Psi_n(p_1, \dots, p_n)$, which satisfies the equivalence

$$\Gamma \cup \Psi_n(\varphi_1, \dots, \varphi_n) \text{ is inconsistent in } \vdash \iff \Gamma \vdash \{\varphi_1, \dots, \varphi_n\}, \quad (1)$$

for all sets of formulas $\Gamma \cup \{\alpha_1, \dots, \alpha_n\}$ [23] (see also [19, 18]). As expected, the only intermediate logic having a classical inconsistency lemma is CPC (with $\Phi_n := \{\neg(p_1 \wedge \dots \wedge p_n)\}$). This is far from the case of bi-intermediate logics. For example, we prove that every member of $\Lambda(\text{bi-LC})$ has a classical inconsistency lemma witnessed by

$$\Phi_n := \{\sim \neg \sim (p_1 \wedge \dots \wedge p_n)\},$$

where $\neg p$ and $\sim p$ are shorthand for $p \rightarrow \perp$ and $\top \leftarrow p$ (see, e.g., [22, Chpt. 4]). Accordingly, logics in $\Lambda(\text{bi-LC})$ exhibit a certain balance between the classical and intuitionistic behavior of negation connectives.

The logic bi-LC is algebraized by the variety bi-GA of *bi-Gödel algebras*, i.e., the class of bi-Heyting algebras which satisfy Gödel's pre-linearity axiom $(p \rightarrow q) \vee (q \rightarrow p)$. This is a semi-simple variety of bi-Heyting algebras, hence it follows from [29] that it has a discriminator term, and therefore has EDPC. Moreover, as this variety is axiomatized (relative to bi-HA) by a \leftarrow -free formula and has a locally finite Heyting algebra reduct [8], it follows from [22, Chpt. 3] that bi-GA enjoys the finite model property.

As for the geometric models of bi-LC, these take the form of *bi-Esakia co-forests*, i.e., bi-Esakia spaces whose underlying posets are disjoint unions of co-trees. In particular, the dual spaces of the simple bi-Gödel algebras are termed *bi-Esakia co-trees*, and as finite bi-Esakia spaces are equipped with the discrete topology, all finite co-trees can be viewed as a bi-Esakia co-trees.

The main contributions of our work can be summarized as follows. We develop a theory of Jankov, subframe and canonical formulas of bi-Gödel algebras. We employ Jankov formulas to obtain a characterization of splittings in $\Lambda(\text{bi-LC})$ and canonical formulas to uniformly axiomatize all the extensions of bi-LC, cf. [2].

Theorem 1. *If $L \in \Lambda(\text{bi-LC})$, then:*

1. *L is a splitting logic iff L is the logic of a finite co-tree;*

2. L is axiomatizable by canonical formulas. Moreover, if L is finitely axiomatized, then L is axiomatizable by finitely many canonical formulas.

We also use Jankov formulas to show that $\Lambda(\text{bi-LC})$ has the cardinality of the continuum. This is achieved by means of the construction of an infinite antichain (with respect to the order of being a bi-Esakia morphic image) of finite co-trees², and contrasts with the case of $\Lambda(\text{LC})$ which is well known to be a chain of order type $(\omega + 1)^\partial$ [8].

Lastly, subframe formulas can be used to describe the fine structure of co-trees, since a bi-Esakia co-tree \mathcal{X} refutes the subframe formula of (the algebraic dual of) a finite co-tree \mathfrak{F} iff \mathcal{X} admits \mathfrak{F} as a subposet. For the present purpose, the interest of subframe formulas is that they help us characterize the locally tabular extensions of bi-LC. This is done in three steps, all relying on the structure of the *finite combs*, a particular class of co-trees depicted in Figure 1.

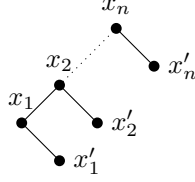


Figure 1: The n -comb \mathfrak{C}_n

Firstly, we prove that for all positive integers n , a bi-Esakia co-tree \mathcal{X} admits \mathfrak{C}_n as a subposet iff \mathfrak{C}_n is a bi-Esakia morphic image of \mathcal{X} . Secondly, we find a natural bound for the size of m -generated simple bi-Gödel algebras whose bi-Esakia duals do not admit the n -comb \mathfrak{C}_n as a subposet. Finally, by showing that the variety generated by the (algebraic duals of the) finite combs is not locally finite, we derive the following criterion and an immediate corollary:

Theorem 2. *If $L \in \Lambda(\text{bi-LC})$, then L locally tabular iff \mathfrak{C}_n is not a model of L , for some $n \in \omega$.*

Corollary 3. *A variety \mathbf{V} of bi-Gödel algebras is locally finite iff \mathbf{V} omits the algebraic dual of a finite comb. Consequently, the variety generated by the duals of the finite combs is the only pre-locally finite variety of bi-Gödel algebras.*

It follows from above that bi-LC is not locally tabular, highlighting yet another contrast with LC, which is well known to be locally tabular [17]. These results are collected in [5].

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